

## Abstract

In this poster, we derive a method of approximating the square root of two. We do this by constructing a geometric figure that has inscribed circles lying across the diagonal of a square. By adding the radii of these circles, we can find the length of the diagonal, which will be the square root of two if the square had side length one. This project uses trigonometry, recursion, the binomial theorem, and a lot of algebra to arrive at the following equation:

$$\sqrt{2} = \lim_{x \rightarrow \infty} \frac{-3 + 6 \sum_{j=0}^x \sum_{i=2j}^{2x+1} \binom{i}{2j} 8^j 3^{i-2j}}{3 + 4 \sum_{j=0}^x \sum_{i=2j+1}^{2x+1} \binom{i}{2j+1} 8^j 3^{i-2j}}$$

We find that the formula gives an approximation that is about 3 decimal places more precise for every iteration of this sequence. While other methods exist for calculating square roots more efficiently, geometric methods such as these are scarcely found and can apply to other problems. Specifically, these methods have potential applications in finding square roots generally, approximating pi, and describing other irrational numbers. Although space filling techniques have been used to find values, this method may be better as it uses circles. Additionally, it should be noted that this function has some properties which may lead to fast calculations using a program with optimizations. Namely, the 8s can be solved using bit manipulation and constructing Pascal's triangle through recursion can eliminate the use of factorials.

## Introduction

Efficient algorithms for computing square roots are essential for various applications, including numerical simulations, graphics rendering, cryptography, and signal processing. As such, there are many methods that have been discovered to approximate these numbers quickly and accurately. We do just this by claiming that the diagonal is approximated by diameters of the circles lying on it as seen in **Figure 2**. By solving for the radii of each circle and adding them up, we approach the length of the diagonal. The following is a demonstration of how this method arrived at the equation listed above.

## Methods

To now find the radii, we noticed that the radius of the original circle is equal to the sum of the vertical magnitudes of the circles' diameter along the diagonal and the final unknown circle's radius, as seen in **Figure 3**. Using some algebra and induction we find the value of the radii of the preceding circles. Combining this with the equation in **Figure 5**, we get **Figure 6**. From here we have the issue of defining the root with itself, but this can be resolved by grouping all terms that have square roots on one side as seen in **Figure 7**. A general equation then arises through categorizing the terms with Pascal's Triangle as is seen in **Figure 8**. Putting this all together we arrive at **Figure 9**.

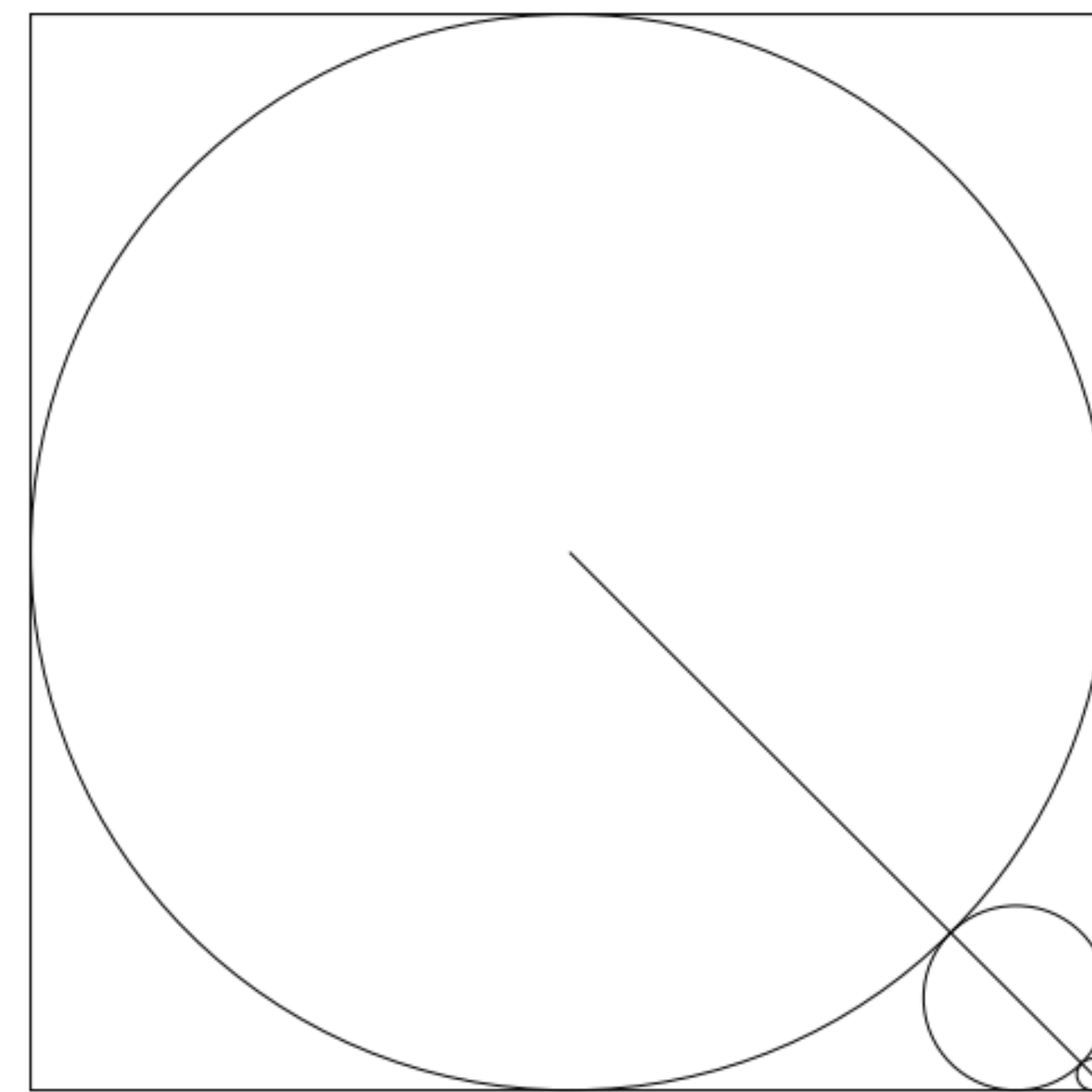


Figure 2. The first three circles lying on the diagonal of a square

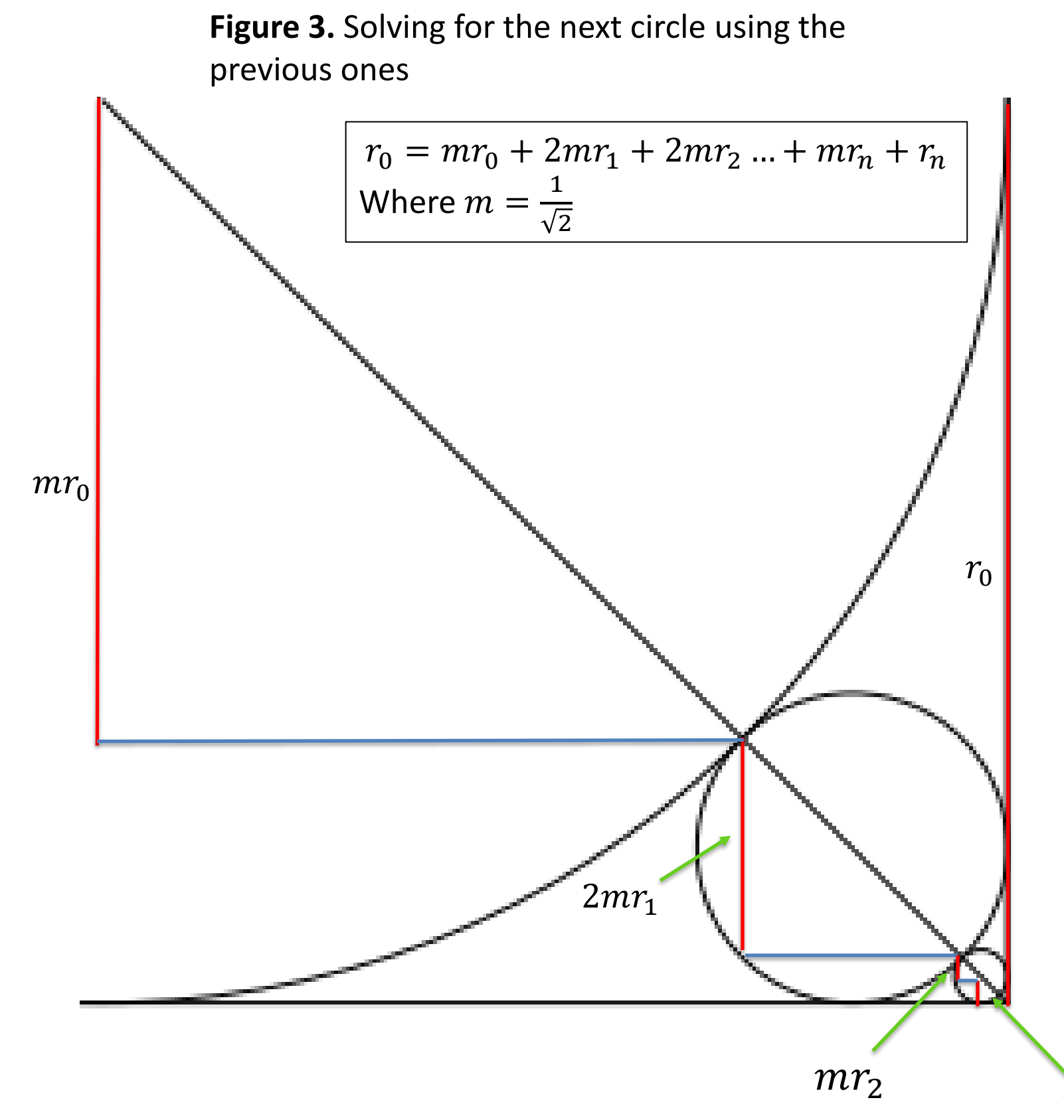


Figure 3. Solving for the next circle using the previous ones

	$r_0 = mr_0 + mr_1 + r_1$
Finding $r_1$	$r_1 = r_0 \frac{1-m}{1+m}$
	$r_0 = mr_0 + 2mr_1 + mr_2 + r_2$
Finding $r_2$	$r_2 = r_0 \frac{1-m-2mr_1}{1+m}$ $r_2 = r_0 \left(\frac{1-m}{1+m}\right)^2$
Generally,	$r_n = r_{n-1} \left(\frac{1-m}{1+m}\right)$
or	$r_n = r_0 \left(\frac{1-m}{1+m}\right)^n$

Figure 4. Solving for the radius

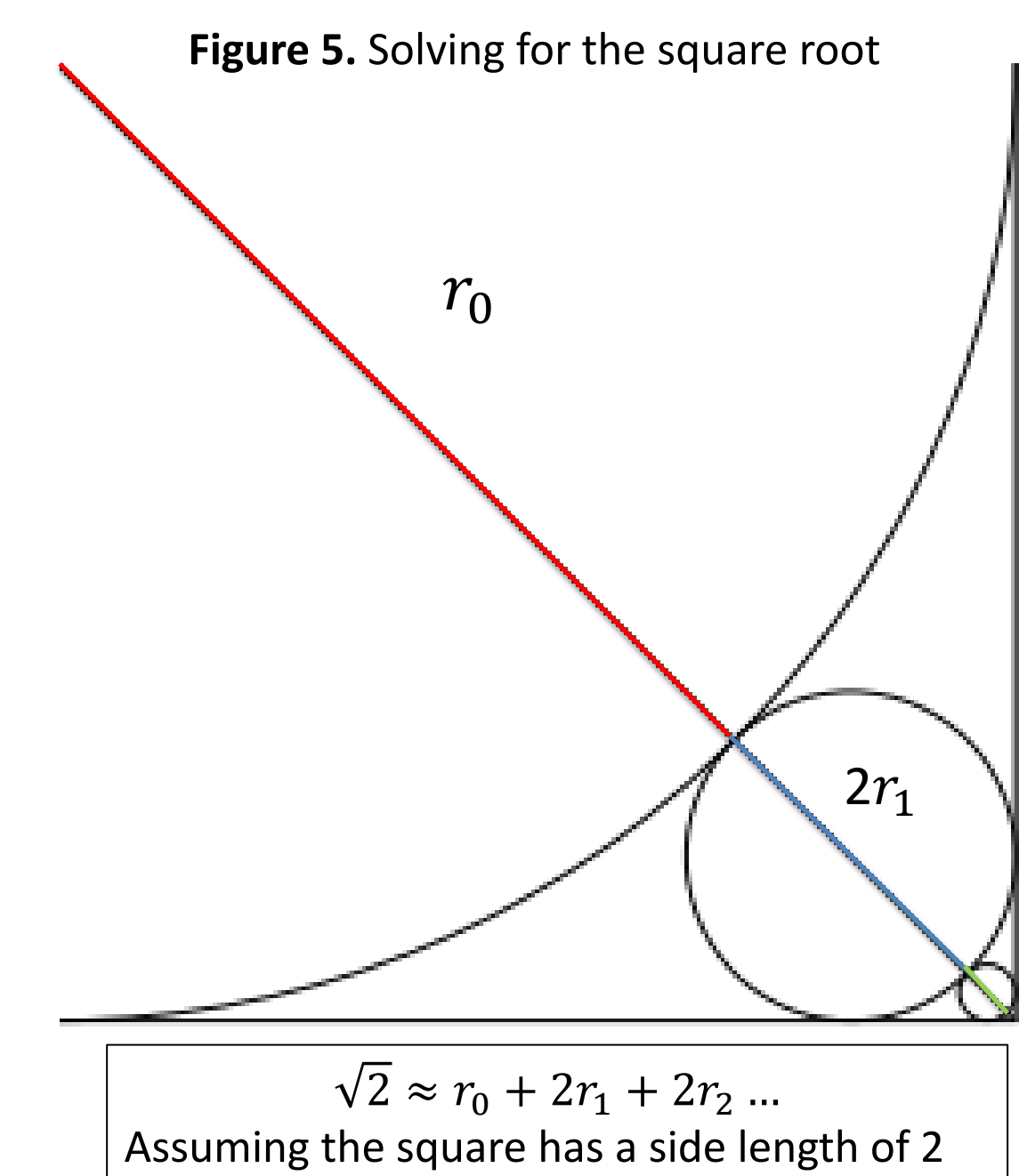


Figure 5. Solving for the square root

$$\sqrt{2} \approx r_0 + 2r_0 \left(\frac{1-m}{1+m}\right) + 2r_0 \left(\frac{1-m}{1+m}\right)^2 + \dots + 2r_0 \left(\frac{1-m}{1+m}\right)^n$$

$$\sqrt{2} \approx \lim_{x \rightarrow \infty} r_0 \left(1 + 2 \sum_{n=1}^x \left(\frac{1-m}{1+m}\right)^n\right)$$

Now because the square has side length 2,  $r_0 = 1$ .

Also  $\frac{1-m}{1+m} = \frac{1-\frac{1}{\sqrt{2}}}{1+\frac{1}{\sqrt{2}}} = 3 - 2\sqrt{2}$

$$\sqrt{2} = \lim_{x \rightarrow \infty} 1 + 2 \sum_{n=1}^x (3 - 2\sqrt{2})^n$$

Figure 6. Combining equations from 5. and 4.

$$\sqrt{2} \approx 1 + 2 \sum_{n=1}^2 (3 - 2\sqrt{2})^n$$

$$\sqrt{2} \approx 1 + 2(3 - 2\sqrt{2} + 9 - 12\sqrt{2} + 8)$$

$$\sqrt{2}(1 + 4 + 24) \approx 1 + 6 + 18 + 16$$

$$\sqrt{2} \approx 1.413793$$

Figure 7. An example of grouping the roots

		1					
	1	1					
	1	2	1				
	1	3	3	1			
	1	4	6	4	1		
	1	5	10	10	5	1	
	1	6	15	20	15	6	1

The terms we are interested in moving are those that have  $-2\sqrt{2}$  as odd power. By grouping the highlighted terms above we get

$$\sum_{j=0}^x \sum_{i=2j+1}^{2x+1} \binom{i}{2j+1} (-2\sqrt{2})^{2j+1} 3^{i-(2j+1)} + \sum_{j=0}^x \sum_{i=2j}^{2x+1} \binom{i}{2j} (-2\sqrt{2})^{2j} 3^{i-2j}$$

Figure 8. Grouping terms using the binomial theorem

Note: -1 is present because the sequence starts from  $n = 0$  instead of  $n = 1$

$$\sqrt{2} = \lim_{x \rightarrow \infty} 1 + 2 \left( -1 + \sum_{j=0}^x \sum_{i=2j+1}^{2x+1} \binom{i}{2j+1} (-2\sqrt{2})^{2j+1} 3^{i-(2j+1)} + \sum_{j=0}^x \sum_{i=2j}^{2x+1} \binom{i}{2j} (-2\sqrt{2})^{2j} 3^{i-2j} \right)$$

$$\sqrt{2} \left( 1 - \frac{4}{3} \lim_{x \rightarrow \infty} \sum_{j=0}^x \sum_{i=2j+1}^{2x+1} \binom{i}{2j+1} 8^j 3^{i-2j} \right) = -1 + 2 \lim_{x \rightarrow \infty} \sum_{j=0}^x \sum_{i=2j}^{2x+1} \binom{i}{2j} 8^j 3^{i-2j}$$

$$\sqrt{2} = \lim_{x \rightarrow \infty} \frac{-3 + 6 \sum_{j=0}^x \sum_{i=2j}^{2x+1} \binom{i}{2j} 8^j 3^{i-2j}}{3 + 4 \sum_{j=0}^x \sum_{i=2j+1}^{2x+1} \binom{i}{2j+1} 8^j 3^{i-2j}}$$

Figure 9. The Final Form

## Results

**Results:** While the final function is written mathematically, it would be far more efficient to go back to **Figure 8** and implement the idea into a computer program. Nonetheless, upon writing the function into Mathematica, some fascinating features were discovered with the algorithm:

Let the function be defined as  $f(x) = \frac{-3 + 6 \sum_{j=0}^x \sum_{i=2j}^{2x+1} \binom{i}{2j} 8^j 3^{i-2j}}{3 + 4 \sum_{j=0}^x \sum_{i=2j+1}^{2x+1} \binom{i}{2j+1} 8^j 3^{i-2j}}$

**1. Error**  
To find how much more precise a function is from term to term you simply find the error from one term then divide it by the error from the next

$$E(x) = \frac{\sqrt{2} - f(x)}{\sqrt{2} - f(x+1)}$$

This function approaches 1154. Upon evaluating  $\log_{10} 1154$  we discover that the function is about 3.062205809 more digits precise every iteration

**2. Actual Values**

While the previous number is just found analytically, it is necessary to check through to see if results look that way. For the value of  $x = 1000$ ,  $f(x)$  produces

891329695472716469343730653727341132532254981841808425782404585654099342713407007624862721396959939810862045132  
349784098205465153770697208376499423680079724346897792926219671955340911876147027508240789572804962813591318164  
2280077984317970606002461738701320455114868811301054963719874030786970112275050605387934136365357589919853458904623646  
448072928237246438675222979142588988845434848963232038347809383792775978450957466599720014021507749461010287089309  
632873040640372408745858600080291624822589157650038923128262788310271633773426744219455762581564408934107568910794469  
475090523765699019672636800915739309837952547906176985188749224825791574110367754300231857483380080870400331672173  
3537025387546978043241761301906074304337346058660742058889790280160152323257390992863347268991660258053397808282  
2454104682930253787867924267878132013039526848482624203857929984418854377481832111430961143139251838160983325  
6182808475206480402001516367334998356550650505281831561492414443753238762113116595400795902433956790802120016  
36820425430525883587352156584223797310804430554389158608300796426881468413910885484879587410035164820000089349  
488155086403165445762050653752023397074186261271861316141775186499220859887011539131265665993510212334842660590331919  
6507636692266619813356626350493052412987758384357319893290259956282168908121161711552267400761338965956673154237663822  
752660570973103671648609345039660142496511827684877208407637963152073508055341089884324463579853614279

as a numerator and

6302652453874561942304247387079941733781357661402512430291056141105072647084442392692534509062996091690330077993742  
165484542419097984724681879407180267398851818441179742181687876578486507065931308467107909610847768382681818638085  
9629709842701135514393098608738648387846236489799417012972181851832604944523876331104038749697596761445280878  
20373798778152723099322889539496506441654368466000205614762754650503584771921101893118418335484871384615813112829281  
77808819279796467647603152520242016919635962977828570801305247352984058393271365773181505369046369883750245783893  
80120318908979891997533684514204295451392998008029289980361352582055674409553162628652964614643492217695387251  
2949475930404711038769917582729571374193931113796160907118553983608167912957083081672473462857870554741719973  
516744434123524885239310792932703606063213354000485403719198931548935514727437539882825593742673967528623091814487  
7560756975294519538575347481377860953419455810198912435274022307091058186907068027641049094653912960250159238202189285  
1812043426847420493765213828812262696968369040724099096959862324827548949675716371451528970145593903423974029952  
165745526715842864722997574070695451541066285505515102008023708463892947705234008803806271298923733727448  
6176969940227352305076955026867914253920715557899154509746765543917375586534691595561818168960520935063427814798373  
958706970546123259420652366660956011760293128518461039655163353566149250635499497511366945775191589

as a denominator; Which produces an error of  $8.900372405 \dots \cdot 10^{-3065}$ . Not bad!

Note: The initial values of this sequence gave approximations that were about 4 more decimal places precise each iteration for some fascinating ones.

## Future Work

1. Generalizing this method to find any square root function.
2. Experiment with various configurations of shapes to approximate transcendental numbers\*.
3. Find a more rapidly converging sequence by using a "tighter" fit such as a rotated ellipse in this square.
4. Calculating the square root of two to millions of digits
5. Expand the method to multiple dimensions (Spheres and Cubes)

\*There is a fair chance this cannot be done as transcendental numbers cannot be constructed with algebraic numbers alone